AI505 Optimization

Constrained Optimization

Marco Chiarandini

Department of Mathematics & Computer Science University of Southern Denmark

Constrained Optimization

- Minimizing an objective subject to design point restrictions called constraints
- A variety of techniques transform constrained optimization problems into unconstrained problems
- New optimization problem statement

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\begin{array}{l} \underset{x}{\text{minimize } f(x)}\\ \text{subject to } x \in \mathcal{X} \end{array}
```

• The set $\mathcal{X} \subset \mathbb{R}$ is called the **feasible set**

Constrained Optimization

Constraints that bound feasible set can change the optimizer



Constraint Types

- Generally, constraints are formulated using two types:
 - 1. Equality constraints: $h(\mathbf{x}) = 0$
 - 2. Inequality constraints: $g(x) \leq 0$
- Any optimization problem can be written as

 $\begin{array}{l} \underset{\boldsymbol{x}}{\text{minimize } f(\boldsymbol{x})}\\ \text{subject to } g_i(\boldsymbol{x}) \leq 0 \text{ for all } i \text{ in } \{1, \ldots, m\}\\ h_j(\boldsymbol{x}) = 0 \text{ for all } j \text{ in } \{1, \ldots, \ell\} \end{array}$

 $\begin{array}{l} \underset{x}{\text{minimize }} f(x) \\ \text{subject to } g(x) \leq 0 \\ h(x) = 0 \end{array}$

f and the functions h and g are all smooth, real-valued functions on a subset of Re^n

Transformations to Remove Constraints

- If necessary, some problems can be reformulated to incorporate constraints into the objective function
- If x is constrained between a and b

$$x = t_{a,b}(\hat{x}) = \frac{b+a}{2} + \frac{b-a}{2} \left(\frac{2\hat{x}}{1+\hat{x}^2}\right)$$



Transformations to Remove Constraints

Example

and a first set of the

 $f(x) = x\sin(x)$ $(f \circ T_{2.6})(\hat{x})$ minimize $x \sin(x)$ x 2 subject to $2 \le x \le 6$ 10 0 0 $^{-2}$ -4-105 -5 0 10 15 0 5 x Ŷ

$$\underset{\hat{x}}{\text{minimize } t_{2,6}(x) \sin(t_{2,6}(x))}$$
$$\underset{\hat{x}}{\text{minimize } \left(4 + 2\left(\frac{2\hat{x}}{1 + \hat{x}^2}\right)\right) \sin\left(4 + 2\left(\frac{2\hat{x}}{1 + \hat{x}^2}\right)\right)$$

 $(\triangle) = (+ (\triangle))$

Transformations to Remove Constraints

Example

minimize
$$f(x)$$

subject to $h(x) = x_1^2 + x_2^2 + \ldots + x_n^2 - 1 = 0$

• Solve for one of the variables to eliminate constraint:

$$x_n = \pm \sqrt{1 - x_1^2 - x_2^2 - \ldots - x_{n-1}^2}$$

• Transformed, unconstrained optimization problem:

minimize
$$\left(\left[x_1, x_2, \dots, x_{n-1}, \pm \sqrt{1 - x_1^2 - x_2^2 - \dots - x_{n-1}^2} \right] \right)$$

Lagrangian Relaxation

- With only equality constraints, critical points (local minima, global minima, or saddle points optimal) where gradient of f and the gradient of h are aligned
- The method of Lagrangian relaxation is used to optimize a function subject to (equality) constraints
- Lagrangian multipliers refer to the variables introduced by the method denoted by λ

1. Form Lagrangian relaxation

 $\begin{array}{l} \underset{x}{\text{minimize }} f(x) \\ \text{subject to } h(x) = 0 \end{array}$

$$\mathcal{L}(\boldsymbol{x},\lambda) = f(\boldsymbol{x}) - \lambda h(\boldsymbol{x})$$

2. Set $\nabla_{\mathbf{x}} \mathcal{L}(\mathbf{x}, \lambda) = 0$ and $\nabla_{\lambda} \mathcal{L}(\mathbf{x}, \lambda) = 0$ to get

$$abla f(\mathbf{x}) = \lambda \nabla h(\mathbf{x}) \qquad h(\mathbf{x}) = 0$$

3. solve for **x** and λ

Example

minimize
$$-\exp\left(-\left(x_1x_2-\frac{3}{2}\right)^2-\left(x_2-\frac{3}{2}\right)^2\right)$$

subject to $x_1-x_2^2=0$

Lagrangian Relaxation

Intuitively, the method of Lagrange multipliers finds the point x^* where the constraint function is orthogonal to the gradient



Lagrangian Relaxation with Inequality Constraints

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\begin{array}{l} \underset{x}{\text{minimize }} f(x) \\ \text{subject to } g(x) \leq 0 \end{array}
```

 If solution lies at the constraint boundary, the constraint is called active, and the Lagrangian condition holds for a non-negative constant μ:

 $\nabla f(\boldsymbol{x}) + \mu \nabla g(\boldsymbol{x}) = 0$

• If the solution lies within the boundary, the constraint is called **inactive**, and the optimal solution simply lies where

 $\nabla f(\mathbf{x}) = 0$

that is, the Lagrangian condition holds with $\mu=0$

Lagrangian Relaxation with Inequality Constraints

 $\begin{array}{l} \underset{x}{\text{minimize }} f(x) \\ \text{subject to } g(x) \leq 0 \end{array}$

 We create the Lagrangian relaxation such that it goes to ∞ outside the feasibility set (g(x) ≤ 0)):

 $\mathcal{L}_{\infty}(\mathbf{x}) = f(\mathbf{x}) + \infty(g(\mathbf{x}) > 0)$

impractical: discontinuous and nondifferentiable.

• Instead, for $\mu > 0$:

 $\mathcal{L}(\mathbf{x}, \mu \geq 0) = f(\mathbf{x}) + \mu g(\mathbf{x})$

 $\mathcal{L}_{\infty}(\mathbf{x}) = \max_{\mu \geq 0} \mathcal{L}(\mathbf{x}, \mu)$

for x infeasible, $\mathcal{L}_{\infty}(x) = \infty$; for x feasible, $\mathcal{L}_{\infty}(x) = f(x)$

• The new optimization problem becomes

 $\underset{\boldsymbol{x}}{\operatorname{minimize}} \underset{\mu \geq 0}{\operatorname{maximize}} \mathcal{L}(\boldsymbol{x}, \mu)$

This is called the primal problem

Necessary Conditions – KKT Conditions

 $\begin{array}{l} \underset{\pmb{x}}{\text{minimize }} f(\pmb{x}) \\ \text{subject to } \pmb{g}(\pmb{x}) \leq 0 \\ \pmb{h}(\pmb{x}) = 0 \end{array}$

Any critical point **x**^{*} must satisfy the Karush-Kuhn-Tucker conditions

1. primal feasibility: $\boldsymbol{g}(\boldsymbol{x}^*) \leq 0$ and $\boldsymbol{h}(\boldsymbol{x}^*) = 0$

2. dual feasibility: penaliztion is towards feasibility $\mu \geq 0$

3. complementary slackness: either μ_i or $g_i(\mathbf{x}^*)$ is zero.

 $\mu_i g_i(\mathbf{x}^*) = 0$, for i = 1, ..., m.

4. stationarity: objective function tanget to each active constraint

$$\nabla f(\mathbf{x}^*) + \sum_i \mu_i \nabla g_i(\mathbf{x}^*) + \sum_j \lambda_j \nabla h_j(\mathbf{x}^*) = 0$$

Necessary Conditions – KKT Conditions

Particular cases

- f concave, g convex: then KKT are also sufficient
- Patological cases

In vector form:

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\left\{egin{aligned} 
abla f(oldsymbol{x}^*)+\mu\cdot
abla oldsymbol{g}(oldsymbol{x}^*)+\lambda\cdot
abla oldsymbol{h}(oldsymbol{x}^*)=0\ oldsymbol{\mu}\cdotoldsymbol{g}(oldsymbol{x}^*)\leq 0,\ oldsymbol{h}(oldsymbol{x}^*)=0\ oldsymbol{\mu}\geq 0 \end{aligned}
ight.
```

Duality

• Generalized Lagrangian Relaxation:

$$\mathcal{L}(\mathbf{x}, \boldsymbol{\mu}, \boldsymbol{\lambda}) = f(\mathbf{x}) + \sum_{i} \mu_{i} g_{i}(\mathbf{x}) + \sum_{j} \lambda_{j} h_{j}(\mathbf{x})$$

• the primal form is

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\underset{\mathbf{x}}{\operatorname{minimize}} \underset{\mu \geq 0, \lambda}{\operatorname{minimize}} \mathcal{L}(\mathbf{x}, \mu. \lambda)
```

• Reversing the order of operations leads to the dual form

 $\underset{\mu \geq 0, \lambda}{\operatorname{maximize minimize }} \mathcal{L}(\boldsymbol{x}, \mu, \lambda)$

• In some cases, the dual problem is easier to solve computationally than the original problem. In other cases, the dual can be used to obtain easily a lower bound on the optimal value of the objective for the primal problem. The dual has also been used to design algorithms for solving the primal problem.

Duality

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Theorem (Max-min inequality)
For any function f: Z \times W \rightarrow \mathbb{R},
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\sup_{z\in Z}\inf_{w\in W}f(z,w)\leq \inf_{w\in W}\sup_{z\in Z}f(z,w).
```

Proof: see wikipedia

- When *f*, *W*, and *Z* are convex the inequality becomes equality and we have a strong max-min property (or a saddle-point property).
- For us:

 $\underset{\mu \geq 0, \lambda}{\operatorname{maximize \ minimize \ } \mathcal{L}(\mathbf{x}, \mu, \lambda) \leq \underset{\mathbf{x}}{\operatorname{minimize \ maximize \ } \mathcal{L}(\mathbf{x}, \mu, \lambda)$

- Therefore, the solution to the dual problem d^* is a lower bound to the primal solution p^*
- The inner part of the dual problem can be used to define the dual function or dual objective $\mathcal{D}(\mu \geq 0, \lambda) = \min_{x} \text{minimize } \mathcal{L}(x, \mu, \lambda)$

Duality

- The dual function is concave. Gradient ascent on a concave function always converges to the global maximum.
- **Dual Problem**: $\max \mathcal{D}(\lambda)$ subject to $\lambda \ge 0$
- Optimizing the dual problem is easy whenever minimizing the Lagrangian with respect to x is easy.
- For any $\mu \geq 0$ and any λ , we have

 $\mathcal{D}(oldsymbol{\mu} \geq 0,oldsymbol{\lambda}) \leq p^*$

- The difference between dual and primal solutions d^* and p^* is called the **duality gap**
- Showing zero-duality gap is a "certificate" of optimality

Penalty methods

• Penalty methods are a way of reformulating a constrained optimization problem as an unconstrained problem by penalizing the objective function value when constraints are violated

Example

 $\begin{array}{l} \underset{\mathbf{x}}{\text{minimize }} f(\mathbf{x}) \\ \text{subject to } \mathbf{g}(\mathbf{x}) \leq 0 \\ \mathbf{h}(\mathbf{x}) = 0 \end{array}$

$$\begin{split} \min_{\mathbf{x}} & f(\mathbf{x}) + \rho \cdot p_{count}(\mathbf{x}) \\ \text{s.t.} & p_{count}(\mathbf{x}) = \sum_{i} (g_i(\mathbf{x}) > 0) + \sum_{j} (h_j(\mathbf{x}) \neq 0) \end{split}$$

Penalty Methods

```
Procedure penalty_method;

Input: f, p, x, k_{max}; \rho = 1, \gamma = 2

Output: x solution

for k in 1, ..., k_{max} do

x \leftarrow minimize_x \{f(x) + \rho \cdot p(x)\};

\rho \leftarrow \rho \cdot \gamma;

if p(x) = 0 then

\lfloor return x;
```

return x;

Penalty methods

• Count penalty:

$$p_{count}(\boldsymbol{x}) = \sum_{i} (g_i(\boldsymbol{x}) > 0) + \sum_{j} (h_j(\boldsymbol{x}) \neq 0)$$

b а ----f(x) $---- f(x) + \rho p_{\text{count}}(x)$ a f(x)

b

a -f(x) $---- f(x) + p_{mixed}(x)$

b $f(x) + \rho p_{\text{quadratic}}(x)$

• Quadratic penalty:

$$p_{quadratic}(oldsymbol{x}) = \sum_i \max(g_i(oldsymbol{x}), 0)^2 + \sum_j h_j(oldsymbol{x})^2$$

• Mixed Penalty:

 $p_{mixed}(\mathbf{x}) = \rho_1 p_{count}(\mathbf{x}) + \rho_2 p_{quadratic}(\mathbf{x})$

Augmented Lagrange Method

• Adaptation of penalty method for equality constraints

$$p_{Lagrangian}(\boldsymbol{x}) \stackrel{def}{=} \frac{1}{2} \rho \sum_{i} (h_i(\boldsymbol{x}))^2 - \sum_{i} \lambda_i h_i(\boldsymbol{x})$$

Procedure augmented lagrange method; Input: $f, h, x, k_{max}; \rho = 1, \gamma = 2$) $\lambda \leftarrow 0;$ for k in $1, \dots, k_{max}$ do $\begin{bmatrix}
p \leftarrow (x \mapsto \rho/2 \cdot \sum_{i} (h_{i}(x)^{2}) - \lambda \cdot h(x)); \\
x \leftarrow \text{minimize}_{x} \{f(x) + p(x)\}; \\
\lambda \leftarrow \lambda - \rho \cdot h(x); \\
\rho \leftarrow \rho \cdot \gamma;
\end{bmatrix}$

return x;

• λ converges towards the Lagrangian multiplier

Interior Point Methods

- Also called barrier methods, interior point methods ensure that each step is feasible
- This allows premature termination to return a nearly optimal, feasible point
- Barrier functions are implemented similar to penalties but must meet the following conditions:
 - 1. Continuous
 - 2. Non-negative
 - 3. Approach infinity as x approaches boundary

Interior Point Methods

• Inverse Barrier:

$$p_{barrier}(\mathbf{x}) = -\sum_{i} \frac{1}{g_i(\mathbf{x})}$$

• Log Barrier:

$$p_{barrier}(m{x}) = -\sum_i egin{cases} \log(-g_i(m{x})) & ext{if } g_i(m{x}) \geq -1 \ 0 & ext{otherwise} \end{cases}$$

New optimization problem:

$$\underset{\mathbf{x}}{\text{minimize } f(\mathbf{x})} + \frac{1}{\rho} p_{\textit{barrier}}(\mathbf{x})$$



Interior Point Methods

Procedure interior point method; Input: $f, \rho, x; \rho = 1, \gamma = 2, \epsilon = 0.001$ $\Delta \leftarrow \infty;$ while $\Delta > \epsilon$ do $x' \leftarrow \text{minimize}_x \{f(x) + \rho(x)/\rho\};$ $\Delta \leftarrow ||x' - x||;$ $x \leftarrow x';$ $\rho \leftarrow \rho \cdot \gamma;$

return x;

- Line searches f(x + αd) are constrained to the interval α = [0, αu], where αu is the step to the nearest boundary.
 In practice, αu is chosen such that x + αd is just inside the boundary to avoid the boundary singularity.
- Needs an initial feasible solutions. Typically, found by solving:

 $\min_{\mathbf{x}} p_{quadratic}(\mathbf{x})$

Summary

- Constraints are requirements on the design points that a solution must satisfy
- Some constraints can be transformed or substituted into the problem to result in an unconstrained optimization problem
- Analytical methods using Lagrange multipliers yield the generalized Lagrangian and the necessary conditions for optimality under constraints
- A constrained optimization problem has a dual problem formulation that is easier to solve and whose solution is a lower bound of the solution to the original problem
- Penalty methods penalize infeasible solutions and often provide gradient information to the optimizer to guide infeasible points toward feasibility
- Interior point methods maintain feasibility but use barrier functions to avoid leaving the feasible set