AI505 Optimization

Discrete Optimization Constraint Programming & Randomized Optimization Heuristics

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2. Randomized Optimization Heuristics

1. Constraint Programming

2. Randomized Optimization Heuristics

Number Circle Puzzle





Put a different number (1 to 8) in each circle such that adjacent circles do not have consecutive numbers.

Example by Patrick Prosser with the help of Toby Walsh, Chris Beck, Barbara Smith, Peter van Beek, Edward Tsang, ...

Which nodes are hardest to number?





Which are the least constraining values to use?



Values 1 and 8



Values 1 and 8



Symmetry means we don't need to consider: 8 1



We can now eliminate many values for other nodes









By symmetry









By symmetry





Value 2 and 7 are left in just one variable domain each









Guess a value, but be prepared to backtrack ...



Guess a value, but be prepared to backtrack ...







Guess another value ...



Guess another value ...







One node has only a single value left ...



Solution



The Core of Constraint Computation

- Modelling
 - Deciding on variables/domains/constraints
- Heuristic Search
- Inference/Propagation
- Symmetry
- Backtracking

Hardness

• The puzzle is actually a hard problem - NP-complete



Constraint programming

- Model problem by specifying constraints on acceptable solutions:
 - define variables and domains
 - post constraints on these variables
- Solve model
 - choose algorithm
 - incremental assignment / backtracking search
 - complete assignments / stochastic search
 - design heuristics

Constraint Satisfaction Problem



- Variable x_i for each node $i = 1, \ldots, 8$
- Domain $\{1, \ldots, 8\}$ for each variable x_i
- Constraints:

$$\begin{split} & \text{allDifferent} \big([x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8] \big) \\ & |x_1 - x_2| > 1 \\ & |x_2 - x_3| > 1 \\ & |x_3 - x_4| > 1 \\ & \vdots \\ & |x_7 - x_8| > 1 \end{split}$$

The **domain** of a variable x, denoted D(x), is a finite set of elements that can be assigned to x.

A constraint *C* on *X* is a subset of the Cartesian product of the domains of the variables in *X*, i.e., $C \subseteq D(x_1) \times \cdots \times D(x_k)$. A tuple $(d_1, \ldots, d_k) \in C$ is called a solution to *C*.

Equivalently, we say that a solution $(d_1, ..., d_k) \in C$ is an assignment of the value d_i to the variable x_i for all $1 \le i \le k$, and that this assignment satisfies C.

If $C = \emptyset$, we say that it is **inconsistent**.

Constraint Satisfaction Problem (CSP)

A CSP is a finite set of variables X with **domain extension** $\mathcal{D} = D(x_1) \times \cdots \times D(x_n)$, together with a finite set of constraints \mathcal{C} , each on a subset of X. A **solution** to a CSP is an assignment of a value $d \in D(x)$ to each $x \in X$, such that all constraints are satisfied simultaneously.

Constraint Optimization Problem (COP)

A COP is a CSP \mathcal{P} defined on the variables x_1, \ldots, x_n , together with an objective function $f: D(x_1) \times \cdots \times D(x_n) \to Q$ that assigns a value to each assignment of values to the variables. An **optimal solution** to a minimization (maximization) COP is a solution d to \mathcal{P} that minimizes (maximizes) the value of f(d).

Another Example: Social Golfers

Example (Social Golfer Problem (Combinatorial Design))

- 9 golfers: 1, 2, 3, 4, 5, 6, 7, 8, 9
- wish to play in groups of 3 players in 4 days
- such that no golfer plays in the same group with any other golfer more than just once.

Is it possible?

	Group 1	Group 2	Group 3
Day 0	???	???	???
Day 1	???	???	???
Day 2	???	???	???
Day 3	???	???	???

This is an instance of a **constrained satisfaction problem**. Adding an optimizing criterion we get a **constrained optimization problem**.

Solution Paradigms

- 1. Dedicated algorithms (eg.: enumeration, branch and bound, dynamic programming)
- 2. Constraint Programming
- 3. Integer Linear Programming
- 4. Other modeling (SAT, SMT, etc.)
- 5. Randomized Search/Optimization Heuristics

Common to 2-5: Representation (modeling) + reasoning (search + inference)

Constraint Programming: Representation

Groups

	Day 0	Day 1	Day 2	Day 3
Golfer 0	1	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 1	1	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 2	1	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 3	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 4	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 5	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 6	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 7	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 8	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}

Golfers

	Group 1	Group 2	Group 3
Day 0	???	???	???
Day 1	???	???	???
Day 2	???	???	???
Day 3	???	???	???

Integer variables:

assign[i,j] variables whose value is from domain $\{1,2,3\}$

Constraints:

- C1: each group has exactly groupSize players
- C2: each pair of players only meets once

Constraint Programming: Representation

```
int: golfers = 9:
int: groupSize = 3;
int: days = 4:
int: groups = golfers/groupSize;
set of int: Golfer = 1..golfers:
set of int: Day = 1..days;
set of int: Group = 1..groups;
array[Golfer, Day] of var Group: assign; % Variables
constraint
 % C1: Each group has exactly groupSize players
 forall (gr in Group, d in Day) ( % c1
   sum (g in Golfer) (bool2int(assign[g,d] = gr)) = groupSize
 ) /\
 % C2: Each pair of players only meets at most once
 forall (g1, g2 in Golfer, d1, d2 in Day where g1 != g2 /\ d1 != d2) (
   (bool2int(assign[q1,d1] = assign[q2,d1]) + bool2int(assign[q1,d2] = assign[q2,d2])) <=1);
solve :: int search([assign[i, j] | i in Golfer, j in Day ],
                   first fail, indomain min, complete) satisfy;
```

Constraint Programming: Reasoning



The solution process proceeds by propagating the constraints on the domanins of the variables (ie, removing values) and tentatively assigning variables until only feasible values are left or backtracking.

1. Constraint Programming

2. Randomized Optimization Heuristics

Yet Another Example: TSP

Example (Traveling Salesman Problem)



Can you find a better solution?

Heuristics

- Get inspired by approach to problem solving in human mind
 [A. Newell and H.A. Simon. "Computer science as empirical inquiry: symbols and search."
 Communications of the ACM, ACM, 1976, 19(3)]
 - effective rules without theoretical support
 - trial and error
- Applications:
 - Optimization
 - But also in Psychology, Economics, Management [Tversky, A.; Kahneman, D. (1974). "Judgment under uncertainty: Heuristics and biases". Science 185]
- Basis on empirical evidence rather than mathematical logic. Getting things done in the given time.

Randomized Optimization Heuristics (ROHs)

Two main search paradigms:

- Constructive search
- Local search

plus high level guiding heuristics (ie, metaheuristics), eg, evolutionary algorithms.



ROHs: Representation

	Group 1	Group 2	Group 3
Day 0	012	345	678
Day 1	0 4 6	1 3 7	258
Day 2	0 4 8	156	2 3 7
Day 3	057	138	246

- Variables = solution representation, tentative solution
- Constraints: relaxed = soft
- Evaluation function to guide the search

ROHs: Reasoning, Local Search Solution: Trial and Error

	Group 1	Group 2	Group 3
Day 0	012	345	678
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Heuristic algorithms: compute, efficiently, **good** solutions to a problem (without caring for theoretical guarantees on running time and approximation quality).

ROHs: Reasoning, Local Search

Example on Traveling Salesman Problem:



ROHs: Reasoning, Metaheuristics







Trying different changes







ROHs: Reasoning, Metaheuristics

- Stochastic Local Ssearch
- Simulated Annealing
- Iterated Local Search
- Tabu Search
- Variable Neighborhood Search
- Adaptive Large Neighborhood Search
- Evolutionary Algorithms
- Ant Colony Optimization
- Estimation-of-Distribution Algorithms
- Artificial Immune Systems
- ...
- Evolutionary Computation Bestiary http://fcampelo.github.io/EC-Bestiary/
- Supernatural inspired [Maturana, Fouhey, 2013]

- White box optimization: models can be expressed mathematically
- Grey box optimization: internal information about objective function computation is often available models that have a mathematical expression but may need data to determine them (eg, neural networks)
- Black box optimization: no mathematical expression is available

Approaches to ROHs

- White/Grey box: representation (modelling) + reasoning (search) constraint based local search, comet, local solver (Hexaly)
- Black Box: a different approach, framework separating problem from solvers and defining the interface specification EasyLocal, ...,

 \implies Cost Action: ROAR-NET



https://github.com/roar-net/roar-net-api-spec

A Search Problem

Definition (Problem statement)

Assume we want to solve a constrained optimization problem: $\min f(x) | x \in F$ where is a set of feasible solutions and f an objective function. All parameters of the problem are known and deterministic.

Definition (Search or Optimization Algorithm)

Goal formulation: we want to find the minimum with respect to some criteria from a set of candidate elements.

Problem formulation: Given a description of the states, an initial state and actions necessary to reach the goal, find a sequence of actions to reach the goal.

Search: the algorithm simulates sequences of actions in the model of the goal, searching until it finds a sequence of actions that reaches the goal. The algorithm might have to simulate multiple tentative answers that do not meet the goal, but eventually it reach a solution, or it will find that no solution is possible.

Search Algorithms

Components of a Search Algorithm (1):

- State or Search Space A set of possible states that the search can be in.
- State or (Candidate) solution: a definition of the states of the search,
- Initial State that the search starts in. For example: an empty set of actions or a complete set of actions.
- Goal A set of one or more goal states. Sometimes there is one goal state sometimes there is a small set of alternative goal states
- Evaluation function f(s) assess the distance from a potential goal. It can also include relaxed constraints.

Search Algorithms

Components of a Search Algorithm (2):

• Action Type *t* available to the algorithm.

Neighborhood Structure

- For a given Action Type t and a State/Solution s, Actions(t, s) returns a finite set of actions of type t that can be executed in s. We say that each of these actions is applicable in s. A transition model, which describes what each action does. Neighborhood
- Result(s, a) returns the state that results from doing action a in state s. Apply Move
- Action-Cost(s, a, s') or c(s, a, s') action cost function gives the numeric cost of applying action a in state s to reach state s'. It reflects the evaluation of the state.

Constraint Handling

In the Constraint Based Local Search community, **constraints** in heuristic methods are handled:

- implicitly in the definition of the search space and of the actions
- as one way constraints
- as soft constraints

ie, relaxed in the evaluation function as objectives with large weights or as lexicographically more important objectives

Application Programming Interface (API)

(Here: not meant as Web API, network-based API, or REST API.) The ROAR-NET API Specification is the definition of an interface or protocol between optimization problems seen as black box and their solvers in order to facilitate understanding, reusing and scaling of solution approaches.

We look for a model which

- ... allows one to use off the shelf components to solve it.
- ... assumes a separation between problem specifics and solver.
- .. is designed as a **software interface** offering a service to other pieces of software and is implemented by the user.
- ... promotes **reusability** of software components and minimizes the user's effort to deploy a solution for the specific optimization problem at hand.
- ... aims at maximizing code extensability, reusability, and simplicity.

Types

- **Problem** the problem instance
- Solution implements the representation of a tentative solution
- Value represents points in objective space
- Neighborhood a function that given a solution, gives another solution neighbor: S → S, based on a neighborhood, compute a Move
- Move applied to a solution to get the novel solution
- **Increment** represents points in objective space

Simplification in single objective cases: **Value** and **Increment** are Reals (ie, **Double** or **Integer**)

Operations: Problem Representation

... on Problem (P):

- empty_solution(P): Solution
- random_solution(P): Solution
- heuristic_solution(P): Solution[0..1]

... on Solution S:

- copy_solution(S): Solution
- lower_bound(S): Value[0..1]
- evaluation_value(): Value[0..1]

Operations: Search

... on Neighborhood (N):

- construction_neighbourhood (Problem) : Neighbourhood
- destruction_neighbourhood(Problem): Neighbourhood
- local_neighbourhood(Problem): Neighbourhood
- moves(N, Solution): Move[0..*]
- random_move(N, Solution): Move[0..1]
- random_move_without_replacement(N, Solution): Move[0..*]

Operations: Search

... on Move:

- lower_bound_increment(Move, Solution): double[0..1]
- objective_value_increment (Move, Solution): double[0..1]
 apply_move (M, Solution): Solution applies the move to a solution, to get a novel solution
- invert_move (M) : Move computes the inverse of a move to revert it

The Full API

Types

Problem Solution Value Neighborhood Move # Operations on Problem
empty_solution
random_solution
heuristic_solution

Operations on Solution
objective_value
lower_bound
copy_solution

Operations on Neighborhood local_neighbourhood construction_neighbourhood destruction_neighbourhood # Operations on Move
moves
random_move
random_moves_without_replacement

lower_bound_increment
objective_value_increment

apply_move
invert_move